

BUS3150 - Computer Facilities Network Management

10 - Internet Protocols

Faculty of Information Technology
Monash University

This Lecture

- Principles of internetworking.
- Internet Protocol Version 4 (IPv4).
- IPv4 Address structure including subnetworking.

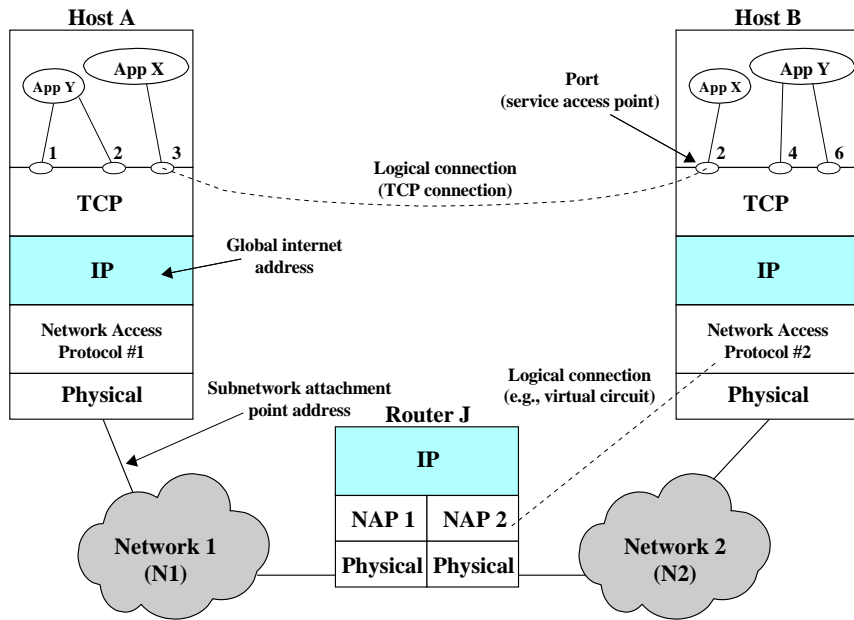
Internetworking

- Techniques to connect various sub-networks so that services can be shared and appear to users as a single large network.
- Establish seamless communication links among machines operating at distant locations, connected to varying types of sub-networks (e.g. LANs) and using different operating systems.
- Require a globally accepted method for connecting networks.

Internetworking

- Communications Network
 - Facility that provides data transfer service.
- An internet
 - Collection of communications networks interconnected by bridges and/or routers.
- The Internet (note upper case I)
 - The global collection of thousands of individual machines and networks.
- Intranet
 - Corporate internet operating within the organization.
 - Uses Internet (TCP/IP and http) technology to deliver documents and resources.

Internetworking Using TCP/IP



Internetworking

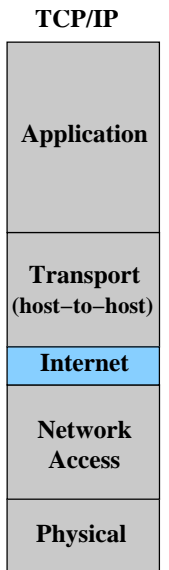
- End System (ES)
 - Device attached to one of the networks of an internet.
 - Supports end-user applications or services.
 - Example: Host A and Host B in the previous slide.
 - Includes workstations, PDA, etc.
- Intermediate System (IS)
 - Device used to connect two networks.
 - Permits communication between end systems attached to different networks.
 - Example: Router J in the previous slide.
 - Includes routers (using IP layer) and bridges (using MAC layer).

Internetworking

- Modes of Operation:
 - Connectionless data transfer
 - Each PDU treated independently.
 - Without connection establishment or termination.
 - Datagram service.
 - Connection-oriented data transfer
 - Involves establishment and termination of a connection.
 - Circuit and virtual circuit connections.
 - Connection-oriented preferred for lengthy exchanges of data.

Internet Protocol (IP) Version 4

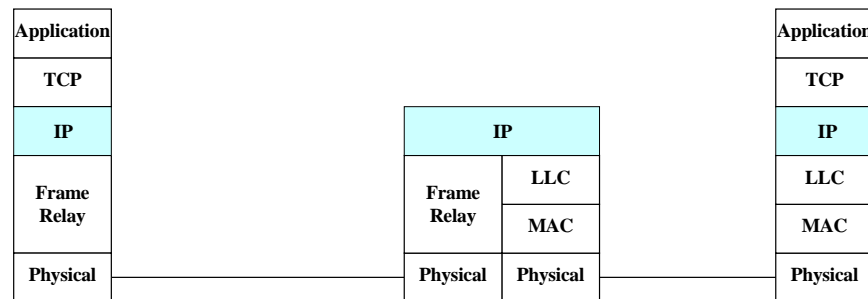
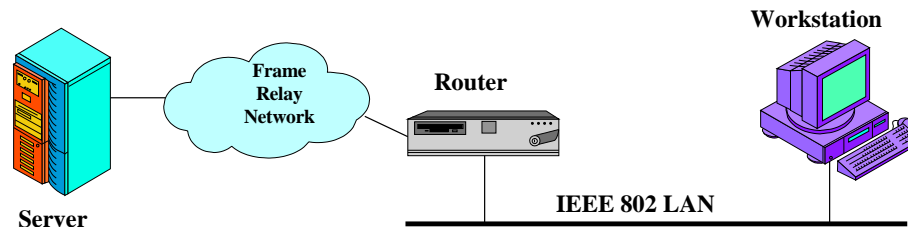
- IPv4 (RFC791) is part of the TCP/IP suite.
- Specifies interface with higher layers (e.g. TCP, UDP).
- Specifies protocol format and mechanisms.
- Will (eventually) be replaced by IPv6.
- Recall the five layers of the TCP/IP suite →
- We are concerned with the Internet layer.



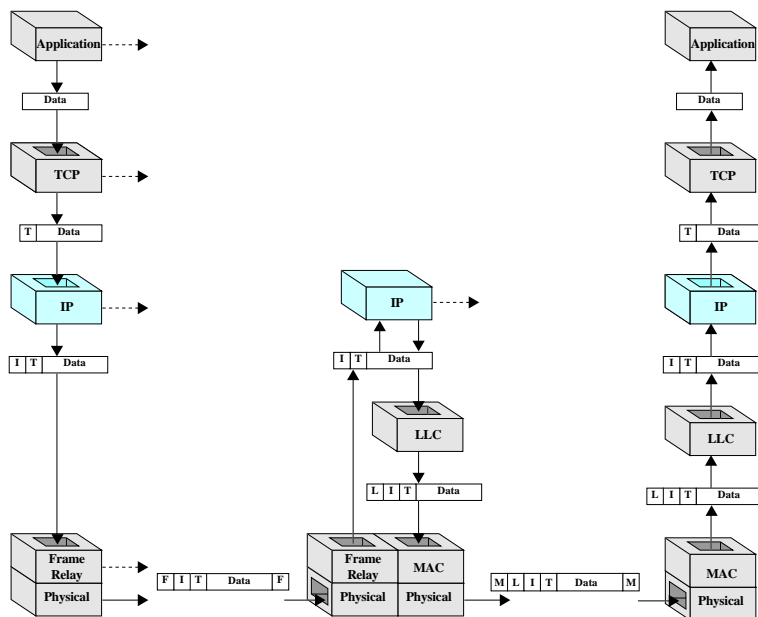
Internet Protocol

- Not all networks within the Internet are reliable.
- The transport layer should deal with reliability issues.
- IP Operation:
 - The transport layer breaks messages into 64kbyte datagrams.
 - Datagrams delivered through Internet routers which may further fragment the datagrams.
 - Uses a best-effort connectionless service to transport datagrams over multiple paths to destination.
 - Transport layer at destination re-orders and re-assembles all datagrams.
- Two protocols in the network layer of each end system and router:
 - Upper layer: internetworking (common over the entire internet).
 - Lower layer: network access (specific to the type of network).

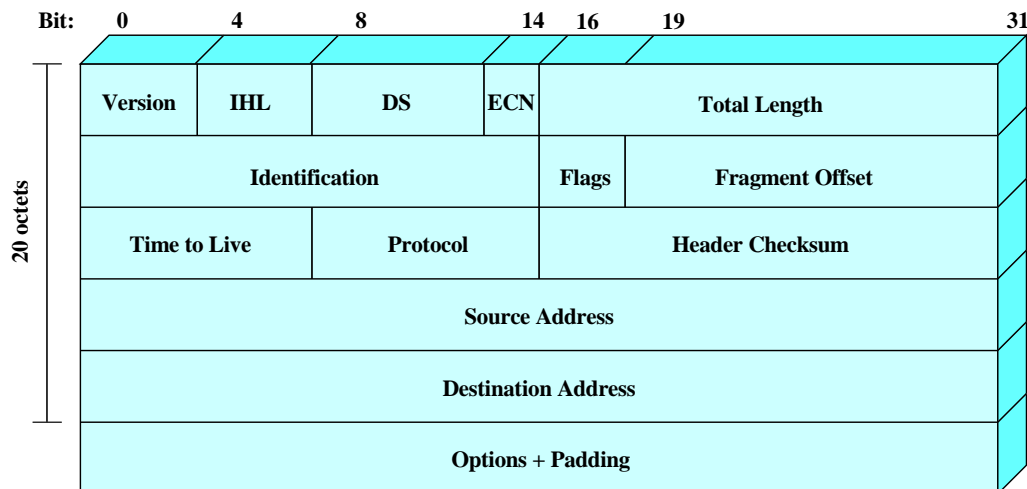
A TCP/IP Example Using Frame Relay and IEEE 802 LAN



A TCP/IP Example Using Frame Relay and IEEE 802 LAN



IPv4 Protocol Header



- IHL : Internet header length.
- DS : Differentiated services field.
- ECN : Explicit congestion notification field.

IP Header Fields

- Version (4 bits)
 - Currently value is 4 for IPv4.
 - For IPv6 the value is 6 and the header is larger (40 octets).
- Internet header length – IHL (4 bits)
 - Multiple of 32 bit words including options.
 - Minimum value of 5 for a minimum header length of 20 octets.
- DS/ECN (8 bits) – precedence, reliability, delay, throughput
 - 6 bits (differentiated services) + 2 congestion notification
- Total length (16 bits)
 - Of the datagram (header and data), in octets.

IP Header Fields

- Identification (16 bits)
 - Sequence number.
 - Used with addresses and user protocol to identify datagram uniquely.
- Flags (3 bits) – Two are only defined.
 - More bit – indicates more datagrams to follow for defragmentation.
 - Don't fragment – If set and datagram is larger than network can handle, then the datagram is dropped. If not set, datagram can be fragmented.
- Fragment offset (13 bits) – multiples of 64 bit words.
 - Each datagram (except the last) should be multiple of 64 bits.
- Time to live (8 bits) – seconds (Each router will decrease time on its way).
- Protocol (8 bits)
 - Next higher layer to receive data field at destination e.g. TCP.

IP Header Fields

- Header checksum (16 bits)
 - Header may change at each node.
 - Re-verified and recomputed at each router.
 - 16 bit ones complement addition of all 16 bit words in header.
 - Set to zero during calculation.
- Source address (32 bits).
- Destination address (32 bits).
- Options – encodes user options (variable length) such as security labels, source routing lists, route recording, etc.
- Padding
 - To ensure the total header length is a multiple of 32 bits.

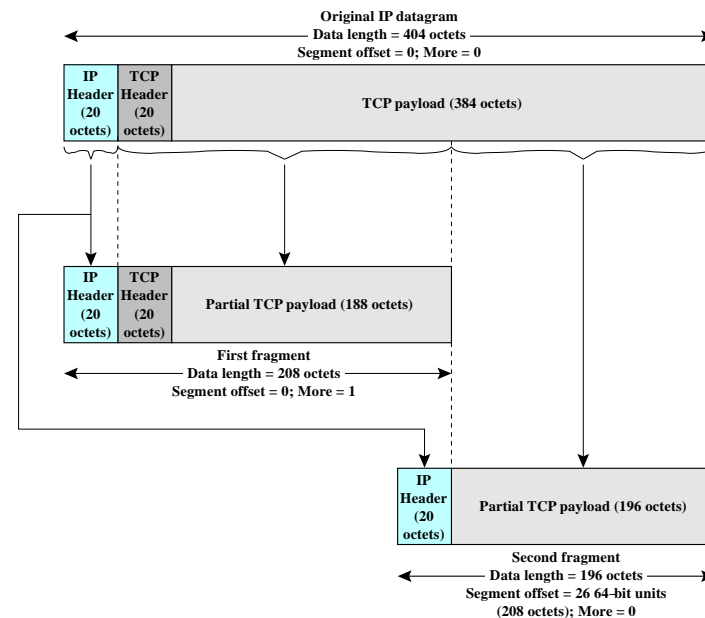
IP Data Field

- Carries user data from next layer up.
- Integer multiple of 8 bits (octet) in length.
- Max length of datagram (header plus data) 65,535 octets.
 - The total length field is 16 bits ($2^{16} - 1 = 65,535$).

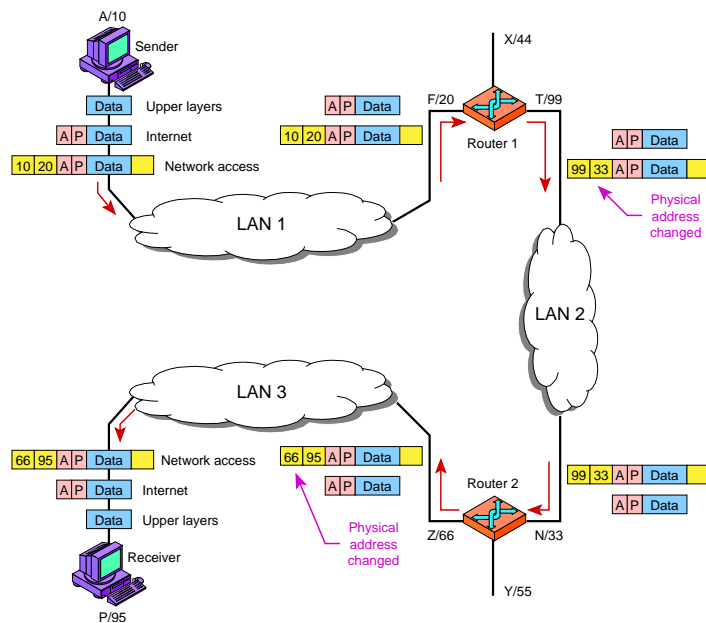
IP Fragmentation

- Different networks have different maximum packet sizes.
- Fragment larger packets – where do we re-assemble fragments?
 - Intermediate systems: need large buffers at each router and require all fragments to follow the same path through the network.
 - Destination: packets get smaller as they traverse the internet.
- IP fragmentation uses re-assembly at destination with:
 - Data unit identifier – use addresses, protocol and identification fields.
 - Data length – length of user data.
 - Fragment offset – position of fragment of user data in original datagram.
 - More flag – indicates this is not the last fragment.

IP Fragmentation Example

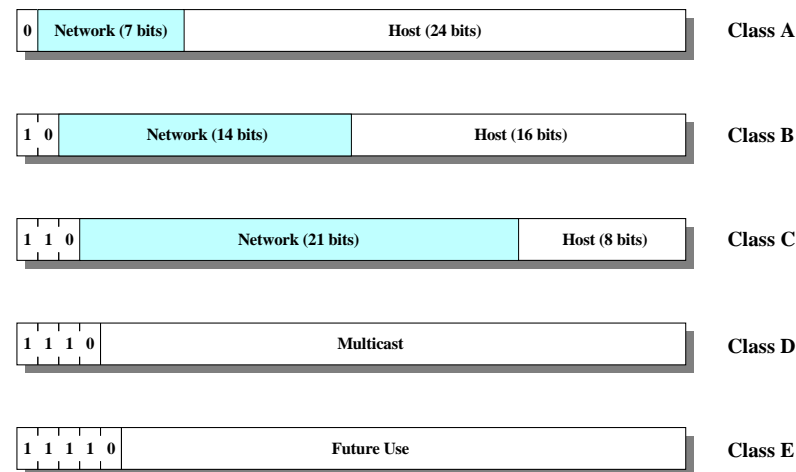


IP Address



IPv4 Address Formats

- 32-bit global address is partitioned into a network (address a subnetwork) and a host (address a host on that subnetwork) identifier.



IPv4 Addresses: Network Class A

- A few networks, each with many hosts.
- Starts with a binary 0.
- Address with all 0 reserved.
- First octet 01111111₂ (127) reserved for loopback.
- Range 1.x.x.x to 126.x.x.x



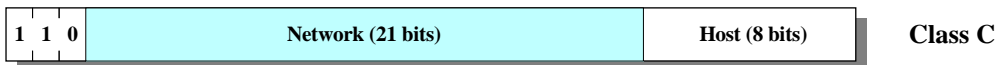
IPv4 Addresses : Network Class B

- Medium number of networks each with a medium number of hosts.
- Starts with a binary 10.
- Range 128.x.x.x to 191.x.x.x
- Second Octet also included in network address:
 - 6 bits from first octet and 8 bits from second octet (6+8=14).
- $2^{14} = 16,384$ class B addresses.
- For example:
 - 130.194.x.x (10000010.11000010.x.x) Monash network.
 - 130.194.11.x Monash ITS managed services.
 - 130.194.11.149 (www.monash.edu.au)



IPv4 Addresses : Network Class C

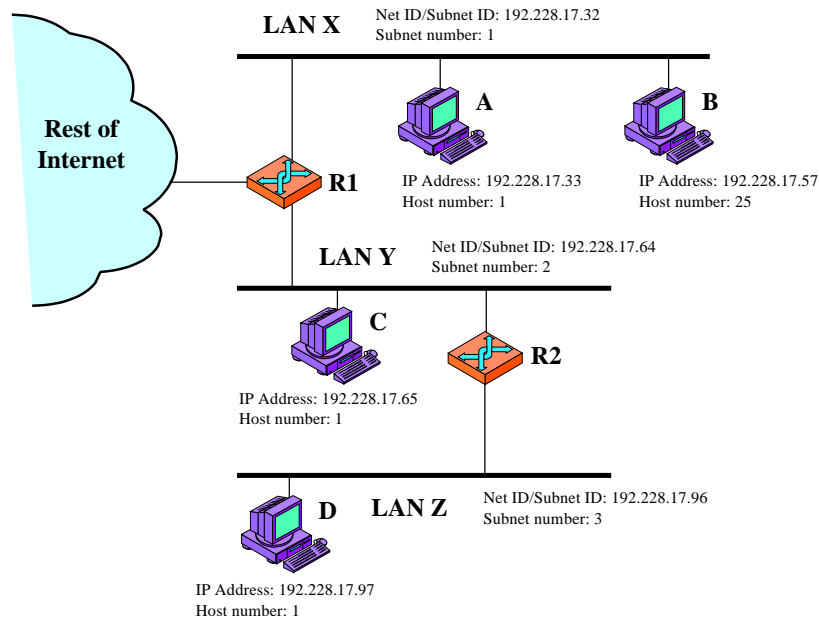
- Many networks, each with a few hosts.
- Starts with a binary 110.
- Range 192.x.x.x to 223.x.x.x
- Second and third octet also part of network address.
- $2^{21} = 2,097,152$ addresses.



Subnets and Subnet Masks

- Allow arbitrary complexity of internetworked LANs within organization.
- Insulate overall internet from growth of network numbers and routing complexity.
- Site looks to rest of internet like single network.
- Each LAN assigned subnet number.
- Host portion of address partitioned into subnet number and host number.
- Local routers route within subnetted network.
- Subnet mask indicates which bits are subnet number and which are host number.

Routing Using Subnetworking



Routing Using Subnetworking

- Uses class C address 192.228.17.x for entire network.
- Upper three bits of host are used to specify the LAN.

	Dotted Decimal	Binary Representation
IP Address	192.228.17.57	11000000.11100100.00010001.00111001
Subnet mask	255.255.255.224	11111111.11111111.11111111.11100000
Bitwise AND of above (network/subnet number)	192.228.17.57	11000000.11100100.00010001.00100000
Subnet number	1 (LAN X)	11000000.11100100.00010001. 001
Host number	25 (Host B)	00000000.00000000.00000000.000 11001

- Subnet mask 255.255.255.0 (class C) delivers to network, while subnet mask 255.255.255.224 delivers to a particular LAN.

Further Reading

- Stallings, W. Data and Computer Communications, Prentice Hall. Chapter 18.
- Stallings, W. Business Data Communications, Prentice Hall. Chapters 4 and 5.
- Forouzan, B. Data Communications and Networking, McGraw-Hill. Parts of Chapters 2, 19 and 20.